# ON GENERAL FIXED POINT METHOD BASED ON MATRIX SPLITTING FOR SOLVING LINEAR COMPLEMENTARITY PROBLEM 

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#### Abstract

In this article, we introduce a modified fixed point method to process the large and sparse linear complementarity problem (LCP) and formulate an equivalent fixed point equation for the LCP and show the equivalence. Also, we provide convergence conditions when the system matrix is a $P$-matrix and two sufficient convergence conditions when the system matrix is an $H_{+}$-matrix. To show the efficiency of our proposed method, we illustrate two numerical examples for different parameters.


MSC. 65F10, 65F50.
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## 1. INTRODUCTION

Given $A_{1} \in \mathbb{R}^{n \times n}$ and a vector $q \in \mathbb{R}^{n}$, the linear complementarity problem denoted as $\operatorname{LCP}\left(q, A_{1}\right)$ is to find the solution $z \in \mathbb{R}^{n}$ to the following system

$$
\begin{equation*}
z \geq 0, A_{1} z+q \geq 0, z^{T}\left(A_{1} z+q\right)=0 \tag{1}
\end{equation*}
$$

The LCP has many applications, including operations research, control theory, mathematical economics, optimization theory, stochastic optimal control, the American option pricing problem, economics, elasticity theory, the free boundary problem, and the Nash equilibrium point of the bimatrix game, which has been extensively studied in the literature on mathematical programming. For details see [14], [10], [21]. For recent works on this problem see [8], [9].

The methods available for solving the LCP may be classified into two groups namely pivoting method [2], [5], [23] and iterative method [22], [18]. The

[^0]basic idea of the pivotal method is to obtain a basic feasible complementary vector through a series of pivot steps, whereas the iterative method generates a series of iterations that converge to a solution [4], [7]. Lemke and Howson [20] introduced the complementary pivot method. Following this method, Lemke introduced a technique known as Lemke's algorithm, which is well known for finding the solution to the LCP.

In order to develop effective iteration methods, we commonly use matrix splittings to find a numerical solution of the large and sparse $\operatorname{LCP}\left(q, A_{1}\right)$, such as the projected methods [16], [19], [25], the modulus algorithm [11] and the modulus based matrix splitting iterative methods [26], [27].

A general fixed point method (GFP) is proposed by Fang [17] assuming the case where $\Omega=\omega D_{1}^{-1}$ with $\omega>0$ and $D_{1}$ is the diagonal matrix of $A_{1}$. The GFP approach takes less iterations than the modulus based successive over relaxation method (MSOR) [11]. We present a modified form of GFP [17] that incorporates projected type iteration techniques by including two positive diagonal parameter matrices $\Omega_{1}, \Omega_{2}$ and $\phi$ is a strictly lower triangular matrix in this article. We also show that the fixed point equation and the LCP is equivalent and discuss the convergence conditions and along with several convergence domains for our method.

The paper is organised as follows: we review some notation, definitions and lemmas in Section 2 in order to establish our key findings. The iterative fixed point approach for solving $\operatorname{LCP}\left(q, A_{1}\right)$ with convergence analysis is proposed in Section 3. We present two numerical examples in Section 4 to demonstrate the efficiency of the proposed methods. Section 5 ends the paper with some conclusions.

## 2. PRELIMINARIES

Some notation, preliminary definitions, and required lemmas are reviewed. Here $A_{1}=\left(\bar{a}_{i j}\right) \in \mathbb{R}^{n \times n}$ and $B_{1}=\left(\bar{b}_{i j}\right) \in \mathbb{R}^{n \times n}$ are square matrices. For $A_{1}=\left(\bar{a}_{i j}\right) \in \mathbb{R}^{n \times n}$ and $B_{1}=\left(\bar{b}_{i j}\right) \in \mathbb{R}^{n \times n}, A_{1} \geq(>) B_{1}$ means $\bar{a}_{i j} \geq(>) \bar{b}_{i j}$ for all $i, j \in\{1,2, \ldots, n\}$.

DEFINITION 1 ([17]). Let $A_{1}=\left(\bar{a}_{i j}\right) \in \mathbb{R}^{n \times n}$ be a square matrix, then $\left|A_{1}\right|=\left(\bar{b}_{i j}\right)$ is defined by $\bar{b}_{i j}=\left|\bar{a}_{i j}\right| \forall i, j$ and $\left|A_{1}\right|$ represent that $\bar{a}_{i j} \geq 0 \forall i, j$.

Definition 2 ([17]). Let $A_{1}, B_{1} \in \mathbb{R}^{n \times n}$ be two square matrices. Then $\left|A_{1}+B_{1}\right| \leq\left|A_{1}\right|+\left|B_{1}\right|$ and $\left|A_{1} B_{1}\right| \leq\left|A_{1}\right| \cdot\left|B_{1}\right|$. Moreover, when $a_{1}, b_{1} \in \mathbb{R}^{n}$ then $\left|a_{1}+b_{1}\right| \leq\left|a_{1}\right|+\left|b_{1}\right|$ and $\left|\left|a_{1}\right|-\left|b_{1}\right|\right| \leq\left|a_{1}-b_{1}\right|$.

Definition 3 ([4]). Let $A_{1} \in \mathbb{R}^{n \times n}$ be a square matrix. $A_{1}$ is said to be a $P$-matrix if all its principle minors are positive such that $\operatorname{det}\left(A_{1 \alpha_{1} \alpha_{1}}\right)>0$ for all $\alpha_{1} \subseteq\{1,2, \ldots, n\}$.

DEfinition 4 ([17]). Suppose $A_{1} \in \mathbb{R}^{n \times n}$ is a square matrix, then its comparison matrix is defined as $\left\langle\bar{a}_{i j}\right\rangle=\left|\bar{a}_{i j}\right|$ if $i=j$ and $\left\langle\bar{a}_{i j}\right\rangle=-\left|\bar{a}_{i j}\right|$ if $i \neq j$.

DEfinition 5 ([6]). Suppose $A_{1} \in \mathbb{R}^{n \times n}$ is a square matrix. It is said to be a Z-matrix if all of its non diagonal elements are less than or equal to zero; an $M$-matrix if $A_{1}^{-1} \geq 0$ as well as $Z$-matrix. The matrix $A_{1}$ is said to be an $H$-matrix if $\left\langle A_{1}\right\rangle$ is an M-matrix. $A_{1}$ is said to be an $H_{+}$-matrix if $A_{1}$ is an $H_{+}$-matrix with $\bar{a}_{i i}>0 \forall i \in\{1,2, \ldots, n\}$.

Definition 6 ([6]). The splitting $A_{1}=M_{1}-N_{1}$ is called an $M$-splitting if $M_{1}$ is a nonsingular $M$-matrix and $N_{1} \geq 0$; an $H$-splitting if $\left\langle M_{1}\right\rangle-\left|N_{1}\right|$ is an $M$-matrix; an $H$-compatible splitting if $\left\langle A_{1}\right\rangle=\left\langle M_{1}\right\rangle-\left|N_{1}\right|$.

Lemma 7 ([1]). Let $a_{1}, b_{1} \in \mathbb{R}^{n}$. Then $a_{1} \geq 0, b_{1} \geq 0, a_{1}^{T} b_{1}=0$ if and only if $a_{1}+b_{1}=\left|a_{1}-b_{1}\right|$.

Lemma 8 ([6]). Suppose $A_{1}, B_{1} \in \mathbb{R}^{n \times n}$. If $A_{1}$ and $B_{1}$ are $M$ and $Z$ matrices respectively with $A_{1} \leq B_{1}$, then $B_{1}$ is an $M$-matrix. If $A_{1}$ is an $H$-matrix, then $\left|A_{1}^{-1}\right| \leq\left\langle A_{1}\right\rangle^{-1}$. If $A_{1} \leq B_{1}$, then $\rho\left(A_{1}\right) \leq \rho\left(B_{1}\right)$.

Lemma 9 ([17]). Let $A_{1} \in \mathbb{R}^{n \times n}$ be an $M$-matrix and $A_{1}=M_{1}-N_{1}$ be an $M$-splitting. Let $\rho$ be the spectral radius, then $\rho\left(M_{1}^{-1} N_{1}\right)<1$.

Lemma 10 ([13]). If splitting is an $H$-compatible of an $H$-matrix, then it is an $H$-splitting but converse is not true.

Lemma 11 ([6]). Suppose $A_{1} \geq 0$. If there exist $v>0 \in \mathbb{R}^{n}$ and a scalar $\alpha_{1}>0$ such that $A_{1} v \leq \alpha_{1} v$, then $\rho\left(A_{1}\right) \leq \alpha_{1}$. Moreover, if $A_{1} v<\alpha_{1} v$, then $\rho\left(A_{1}\right)<\alpha_{1}$.

## 3. MAIN RESULTS

For a given vector $s \in \mathbb{R}^{n}$, we consider the vectors $s_{+}=\max \{0, s\}, s_{-}=$ $\max \{0,-s\}$ and $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right)$, where $\phi$ is a strictly lower triangular matrix, $U_{1}$ is a strictly upper triangular matrix of $A_{1} . U_{1}^{T}$ denotes the transpose of $U_{1}, L_{1}$ is strictly lower triangular matrix of $A_{1}$ and $\alpha$ is a positive real number. In the following theorem we convert the $\operatorname{LCP}\left(q, A_{1}\right)$ into a fixed point equation.

Theorem 12. Let $A \in \mathbb{R}^{n \times n}$ with the splitting $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right)$. If $z=\Omega_{1} s_{+}$and $\omega=\Omega_{2} s_{-}$, then the equivalent formulation of the $\operatorname{LCP}\left(q, A_{1}\right)$ in the form of fixed point equation is

$$
\begin{equation*}
s=\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}-\Omega_{2}^{-1} q \tag{2}
\end{equation*}
$$

Proof. Let $z=\Omega_{1} s_{+}$and $w=\Omega_{2} s_{-}$, and $s=s_{+}-s_{-} . \operatorname{From} \operatorname{LCP}\left(q, A_{1}\right)$

$$
\begin{aligned}
\Omega_{2} s_{-} & =A_{1} \Omega_{1} s_{+}+q \\
s & =s_{+}-\Omega_{2}^{-1}\left(A_{1} \Omega_{1} s_{+}+q\right) \\
s & =\left(I_{1}-\Omega_{2}^{-1} A_{1} \Omega_{1}\right) s_{+}-\Omega_{2}^{-1} q \\
s & =\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}-\Omega_{2}^{-1} q .
\end{aligned}
$$

Based on (2) we propose the following iteration method which is referred to as modified general fixed point method (MGFP) for solving the $\operatorname{LCP}\left(q, A_{1}\right)$, (3) $s^{(k+1)}=\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{(k)}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{(k+1)}-\Omega_{2}^{-1} q$.

Let Res be the Euclidean norm of the error vector, which is defined [17] as follows,

$$
\operatorname{Res}\left(z^{(k)}\right)=\left\|\min \left(z^{(k)}, A_{1} z^{(k)}+q\right)\right\|_{2} .
$$

Consider the nonnegative initial vector $z^{(0)} \in \mathbb{R}^{n}$. The iteration process continues until the iteration sequence $\left\{z^{(k)}\right\}_{k=0}^{+\infty} \subset \mathbb{R}^{n}$ converges. The iteration process stop if $\operatorname{Res}\left(z^{(k)}\right)<10^{-5}$. For computing $z^{(k+1)} \in \mathbb{R}^{n}$ we use the following algorithm.

```
Algorithm 1 (Modified General Fixed Point Method)
    - Given any initial vector \(s^{(0)} \in \mathbb{R}^{n}, \epsilon>0\) and set \(k=0\).
    - for \(k=0,1,2, \ldots\) do
    - \(s_{+}^{(k)}=\max \left\{0, s^{(k)}\right\}\)
    - compute Res \(=\operatorname{norm}\left(\min \left(s_{+}^{(k)}, A_{1} s_{+}^{(k)}+q\right)\right)\)
    - if Res \(<\epsilon\) then
    - \(s=s^{(k)}\)
    - break
    - else
    - Using the following scheme, create the sequence \(s^{(k)}\) :
        \(s_{1}^{(k+1)}=\left(\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{(k)}-\Omega_{2}^{-1} q\right)_{1}\),
    - for \(i=2,3 \ldots, n\) do
        \(\begin{aligned} & s_{i}^{(k+1)}=\left(\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{(k)}-\Omega_{2}^{-1} q\right)_{i} \\ &+\left(\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{(k+1)}\right)_{i} \\ & \text { and set } z^{(k+1)}=\Omega_{1} s_{+}^{(k+1)} .\end{aligned}\)
    - end for
    - end if
    - end for.
```

Moreover, the MGFP provides a general structure for solving $\operatorname{LCP}\left(q, A_{1}\right)$. Using some particular values of the parameter matrices $\Omega_{1}, \Omega_{2}$ and we obtain an iterative method. In particular,

- When $\Omega_{1}=I, \Omega_{2}=\Omega^{-1}$ and $\phi=0$, from (3) we have,

$$
\begin{equation*}
s^{(k+1)}=\left(I_{1}-\Omega\left(D_{1}-U_{1}\right)\right) s_{+}^{(k)}+\Omega L_{1} s_{+}^{(k+1)}-\Omega q, \tag{4}
\end{equation*}
$$

this is a GFP [17].
Theorem 13. Suppose $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right) \in \mathbb{R}^{n \times n}$ and $q \in$ $\mathbb{R}^{n}$. Then $s^{*}$ is a solution of (2) if and only if $z^{*}=\Omega_{1} s_{+}^{*}$ is a solution of $L C P\left(q, A_{1}\right)$.

Proof. Let $s^{*}$ be a solution of (2). Then

$$
\begin{aligned}
s^{*} & =\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{*}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{*}-\Omega_{2}^{-1} q, \\
s^{*} & =\left(I_{1}-\Omega_{2}^{-1} A_{1} \Omega_{1}\right) s_{+}^{*}-\Omega_{2}^{-1} q, \\
\Omega_{2} s_{-}^{*} & =A_{1} \Omega_{1} s_{+}^{*}+q .
\end{aligned}
$$

Since $\Omega_{2} s_{-}^{*} \geq 0$,

$$
A_{1} \Omega_{1} s_{+}^{*}+q \geq 0
$$

Moreover,

$$
\left(\Omega_{1} s_{+}^{*}\right)^{T}\left(A_{1} \Omega_{1} s_{+}^{*}+q\right)=\left(\Omega_{1} s_{+}^{*}\right)^{T}\left(\Omega_{2} s_{-}^{*}\right)=0
$$

and $\Omega_{1} s_{+}^{*} \geq 0$. Therefore $z^{*}=\Omega_{1} s_{+}^{*}$ is a solution of $\operatorname{LCP}\left(q, A_{1}\right)$.
Let $z^{*}=\Omega_{1} s_{+}^{*}, w^{*}=\Omega_{2} s_{-}^{*}$ and $s^{*}=s_{+}^{*}-s_{-}^{*}$. From $\operatorname{LCP}\left(q, A_{1}\right)$

$$
\begin{aligned}
\Omega_{2} s_{-}^{*} & =A_{1} \Omega_{1} s_{+}^{*}+q \\
s^{*} & =s_{+}^{*}-\Omega_{2}^{-1}\left(A_{1} \Omega_{1} s_{+}^{*}+q\right) \\
s^{*} & =\left(I_{1}-\Omega_{2}^{-1} A_{1} \Omega_{1}\right) s_{+}^{*}-\Omega_{2}^{-1} q \\
s^{*} & =\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{*}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{*}-\Omega_{2}^{-1} q .
\end{aligned}
$$

Thus, $s^{*}$ is a solution of (2).
In the following theorem, we show that the solution of (2) is unique when the system matrix $A_{1}$ of $\mathrm{LCP}\left(q, A_{1}\right)$ is a $P$-matrix.

Theorem 14. Let $A_{1}$ be a $P$-matrix and $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right) \in$ $\mathbb{R}^{n \times n}$ and $q \in \mathbb{R}^{n}$. Then for any positive diagonal matrices $\Omega_{1}$ and $\Omega_{2}$, (2) has a unique solution.

Proof. Since $A_{1}$ is a $P$-matrix, for any $q \in \mathbb{R}^{n} \operatorname{LCP}\left(q, A_{1}\right)$ has a unique solution. Let $y^{*}$ and $u^{*}$ be the solutions of (2). Then

$$
\begin{aligned}
s^{*} & =\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{*}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{*}-\Omega_{2}^{-1} q \\
u^{*} & =\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) u_{+}^{*}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} u_{+}^{*}-\Omega_{2}^{-1} q
\end{aligned}
$$

Since $\Omega_{1} y_{+}^{*}=\Omega_{1} u_{+}^{*} \Longrightarrow y_{+}^{*}=u_{+}^{*}$, therefore

$$
y^{*}=u^{*}
$$

In the following, we prove the convergence conditions when $A_{1}$ is a $P$-matrix.
Theorem 15. Let $A_{1}$ be a P-matrix with $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right) \in$ $\mathbb{R}^{n \times n}$ and $q \in \mathbb{R}^{n}$. Assume

$$
\rho\left(\left(I-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)^{-1}\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right|\right)<1
$$

and $s^{*}$ is the solution of (2). Then the sequence $\left\{z^{(k)}\right\}_{k=1}^{+\infty}$ generated by Algorithm 1 converges to $z^{*}$ for any initial vector $s^{(0)} \in \mathbb{R}^{n}$.

Proof. Suppose $A_{1}$ is a $P$-matrix, then $s^{*}$ is a unique solution of (2). Thus

$$
s^{*}=\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right) s_{+}^{*}+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1} s_{+}^{*}-\Omega_{2}^{-1} q .
$$

From (3), this implies

$$
\begin{aligned}
s^{(k+1)}-s^{*}= & \left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\left(s_{+}^{(k)}-s_{+}^{*}\right) \\
& +\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\left(s_{+}^{(k+1)}-s_{+}^{*}\right) .
\end{aligned}
$$

It follows that

$$
\begin{aligned}
& \left|s^{(k+1)}-s^{*}\right|= \\
& \quad=\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\left(s_{+}^{(k)}-s_{+}^{*}\right)+\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\left(s_{+}^{(k+1)}-s_{+}^{*}\right)\right| \\
& \quad \leq\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\left(s_{+}^{(k)}-s_{+}^{*}\right)\right|+\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\left(s_{+}^{(k+1)}-s_{+}^{*}\right)\right| \\
& \quad \leq\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\left|s^{(k)}-s^{*}\right|+\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\left|s^{(k+1)}-s^{*}\right| \\
& \left|s^{(k+1)}-s^{*}\right|-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\left|s^{(k+1)}-s^{*}\right| \leq\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\left|s^{(k)}-s^{*}\right| \\
& \left(I-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)\left|s^{(k+1)}-s^{*}\right| \leq\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\left|s^{(k)}-s^{*}\right|
\end{aligned}
$$

and
$\left|s^{(k+1)}-s^{*}\right| \leq\left(I-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)^{-1}\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\left|s^{(k)}-s^{*}\right|$.
Therefore, if $\rho\left(\left(I-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)^{-1}\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\right)<1$, for any initial vector $s^{(0)} \in \mathbb{R}^{n}$ the sequence $\left\{z^{(k)}\right\}_{k=1}^{+\infty}$ converges to the $z^{*}$.

Now, when $A_{1}$ is an $H_{+}$-matrix, we analyze the convergence domains for parameter matrices $\Omega_{1}$ and $\Omega_{2}$ for MGFP.

Theorem 16. Let $A_{1}$ be a $H_{+}$-matrix with $A_{1}=\left(D_{1}+\phi\right)-\left(L_{1}+U_{1}+\phi\right) \in$ $\mathbb{R}^{n \times n}$ and either one of the following is true:
(1) $\Omega_{2}^{-1} \Omega_{1}>\left(D_{1}+\phi\right)^{-1}$ and $\left(2 \Omega_{2}^{-1} \Omega_{1}-\left(D_{1}+\phi\right)-|B+\phi|\right)$, where $B=$ $L_{1}+U_{1}$.
(2) $0<\Omega_{2}^{-1} \Omega_{1} \leq\left(D_{1}+\phi\right)^{-1}$.

Then the sequence $\left\{z^{(k)}\right\}_{k=1}^{+\infty}$ generated by Algorithm 1 converges to $z^{*}$ for any initial vector $s^{(0)} \in \mathbb{R}^{n}$.

Proof. Since $A_{1}$ is an $H_{+}$-matrix. there $\operatorname{LCP}\left(q, A_{1}\right)$ has unique solution [11]. Now we will look at the splitting,

$$
\begin{aligned}
& \left(I_{1}-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right|= \\
& =\left(I_{1}-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi\right) \Omega_{1}\right|\right)-\Omega_{2}^{-1}|B+\phi| \Omega_{1} .
\end{aligned}
$$

(1) If $\Omega_{2}^{-1} \Omega_{1}>\left(D_{1}+\phi\right)^{-1}$ then,

$$
\begin{aligned}
& \left(I_{1}-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi\right) \Omega_{1}\right|\right)-\Omega_{2}^{-1}|B+\phi| \Omega_{1}= \\
& =2 I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi\right) \Omega_{1}-\Omega_{2}^{-1}|B+\phi| \Omega_{1} \\
& =\Omega_{2}^{-1}\left(2 \Omega_{2}^{-1} \Omega_{1}-\left(D_{1}+\phi\right)-|B+\phi|\right) \Omega_{1} .
\end{aligned}
$$

Since $\left(2 \Omega_{2}^{-1} \Omega_{1}-\left(D_{1}+\phi\right)-|B+\phi|\right)$ is an $M$-matrix. Then the splitting $\left(I_{1}-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right|$ represent an $M$-splitting of the $M$-matrix $\Omega_{2}^{-1}\left(2 \Omega_{2}^{-1} \Omega_{1}-\left(D_{1}+\phi\right)-|B+\phi|\right) \Omega_{1}$, hence $\rho\left(\left(I-\mid \Omega_{2}^{-1}\left(L_{1}+\right.\right.\right.$ $\left.\left.\phi) \Omega_{1} \mid\right)^{-1}\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\right)<1$.
(2) If $\Omega_{2}^{-1} \Omega_{1} \leq\left(D_{1}+\phi\right)^{-1}$ then,

$$
\begin{aligned}
\left(I_{1}-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi\right) \Omega_{1}\right|\right)-\Omega_{2}^{-1}|B+\phi| \Omega_{1} & =\Omega_{2}^{-1}\left(D_{1}+\phi-|B+\phi|\right) \Omega_{1} \\
& =\Omega_{2}^{-1}\left\langle A_{1}\right\rangle \Omega_{1} .
\end{aligned}
$$

Therefore, $\left(I_{1}-\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)-\left|I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right|$ represents an $M$-splitting of $M$-matrix $\Omega_{2}^{-1}\left\langle A_{1}\right\rangle \Omega_{1}$ [13]. Therefore, from Lemma $8 \rho((I-$ $\left.\left.\left|\Omega_{2}^{-1}\left(L_{1}+\phi\right) \Omega_{1}\right|\right)^{-1}\left|\left(I_{1}-\Omega_{2}^{-1}\left(D_{1}+\phi-U_{1}\right) \Omega_{1}\right)\right|\right)<1$.

## 4. NUMERICAL EXAMPLES

In this section, two numerical examples are provided to show the effectiveness of our new method. We use some notation as, $\mathrm{IT}=$ number of iteration steps, CPU $=$ CPU time in seconds. The system matrix $A_{1}$ is generated by

$$
A_{1}\left(p_{1}, p_{2}, p_{3}\right)=Q+p_{1} I_{1}+p_{2} G+p_{3} H,
$$

where $p_{1}, p_{2}$ and $p_{3}$ are given constants, $I_{1}$ is the identity matrix of order $n$ and $G=\operatorname{tridiag}(0,0,1)=\left[\begin{array}{ccccc}0 & 1 & 0 & \ldots & 0 \\ 0 & 0 & 1 & \ldots & 0 \\ 0 & 0 & 0 & 1 & \vdots \\ \vdots & \ldots & 0 & \ddots & 1 \\ 0 & \ldots & 0 & 0 & 0\end{array}\right] \in \mathbb{R}^{n \times n}$ and
$H=\operatorname{diag}([1,2,1,2, \ldots])=\left[\begin{array}{ccccc}1 & 0 & 0 & \ldots & 0 \\ 0 & 2 & 0 & \ldots & 0 \\ 0 & 0 & 1 & 0 & 0 \\ \vdots & \ldots & 0 & 2 & \vdots \\ 0 & \ldots & \ldots & 0 & \ddots\end{array}\right] \in \mathbb{R}^{n \times n}$.
Let $s^{(0)}=(0,0,0,0, \ldots 0,0, \ldots)^{T} \in \mathbb{R}^{n}$ be an initial vector. We consider the $\operatorname{LCP}\left(q, A_{1}\right)$ which has always a unique solution, where $q=(1-1$ $1-1 \ldots 1-1 \ldots)^{T} \in \mathbb{R}^{n}$. We set $\Omega_{1}=I_{1}$ and $\Omega_{2}=\omega^{-1} D_{1}$ in the MGFP. The suggested method is compared with GFP [17], which is effective in solving $\operatorname{LCP}\left(q, A_{1}\right)$.

Matlab version 2021a was used for all calculations. Table 1 and Table 2 show the numerical results for GFP [17] and MGFP respectively.

Example 17. The system matrix $A_{1}, B_{1} \in \mathbb{R}^{n \times n}$ are generated by $A_{1}\left(p_{1}, p_{2}, p_{3}\right)=Q+p_{1} I_{1}+p_{2} G+p_{3} H$, where $p_{1}, p_{2}$ and $p_{3}$ are given constants and

$$
\begin{aligned}
& Q=\operatorname{tridiag}\left(-I_{2}, L_{1},-I_{2}\right)=\left[\begin{array}{ccccc}
L_{1} & -I_{2} & 0 & \ldots & 0 \\
-I_{2} & L_{1} & -I_{2} & \ldots & 0 \\
\vdots & -I_{2} & L_{1} & -I_{2} & \vdots \\
0 & \ldots & -I_{2} & \ddots & -I_{2} \\
0 & \ldots & 0 & -I_{2} & L_{1}
\end{array}\right] \in \mathbb{R}^{n \times n}, \\
& L_{1}=\operatorname{tridiag}(-1,4,-1)=\left[\begin{array}{ccccc}
4 & -1 & \ldots & \ldots & 0 \\
-1 & 4 & -1 & \ldots & 0 \\
\vdots & -1 & 4 & -1 & \vdots \\
0 & \ldots & -1 & \ddots & -1 \\
0 & \ldots & \ldots & -1 & 4
\end{array}\right] \in \mathbb{R}^{m \times m},
\end{aligned}
$$

where $I_{2}$ is the identity matrix of order $m$, where $n=m^{2}$ with $m$ being a positive integer.

Example 18. The system matrix $A_{1} \in \mathbb{R}^{n \times n}$ is generated by

$$
A_{1}\left(p_{1}, p_{2}, p_{3}\right)=Q+p_{1} I_{1}+p_{2} G+p_{3} H,
$$

where $p_{1}, p_{2}$ and $p_{3}$ are given constants, $I_{1}$ is the identity matrix of order $n$ and

$$
\begin{aligned}
& Q=\operatorname{tridiag}\left(-1.5 I_{2}, L_{1},-0.5 I_{2}\right) \\
&=\left[\begin{array}{ccccc}
L_{1} & -0.5 I_{2} & 0 & \ldots & 0 \\
-1.5 I_{2} & L_{1} & -0.5 I_{2} & \ldots & 0 \\
\vdots & -1.5 I_{2} & L_{1} & -0.5 I_{2} & \vdots \\
0 & \ldots & -1.5 I_{2} & \ddots & -0.5 I_{2} \\
0 & \cdots & 0 & -1.5 I_{2} & L_{1}
\end{array}\right] \in \mathbb{R}^{n \times n}, \\
& L_{1}=\operatorname{tridiag}(-1.5,4,-0.5)=\left[\begin{array}{ccccc}
4 & -0.5 & \ldots & \ldots & 0 \\
-1.5 & 4 & -0.5 & \ldots & 0 \\
\vdots & -1.5 & 4 & -0.5 & \vdots \\
0 & \ldots & -1.5 & \ddots & -0.5 \\
0 & \cdots & \cdots & -1.5 & 4
\end{array}\right] \in \mathbb{R}^{m \times m},
\end{aligned}
$$

where $I_{2}$ is the identity matrix of order $m$.
From Table 1 and Table 2, we see that our proposed MGFP have requires less iteration steps than the GFP [17] respectively.

Table 1. Results for GFP [17] and MGFP with $\phi=\alpha\left(L_{1}+U_{1}^{T}\right)$.

| $\mathbf{A}_{\mathbf{1}}(\mathbf{1}, \mathbf{1}, \mathbf{- 1})$ | n | 100 | 400 | 900 | 1600 | 2500 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| GFP | IT | 18 | 21 | 22 | 23 | 23 |
| $\omega=1$ | CPU | 0.0058 | 0.1267 | 0.8356 | 12.8507 | 76.8588 |
|  | Res | $7.8 \mathrm{e}-06$ | $7.5 \mathrm{e}-06$ | $7.3 \mathrm{e}-06$ | $5.3 \mathrm{e}-06$ | $7.0 \mathrm{e}-06$ |
| $\mathbf{M G F P}$ | IT | 15 | 18 | 19 | 20 | 20 |
| $\omega=1$ | CPU | 0.0050 | 0.1028 | 0.6105 | 11.2363 | 67.0586 |
| $\alpha=0.1$ | Res | $6.4 \mathrm{e}-06$ | $7.8 \mathrm{e}-06$ | $6.8 \mathrm{e}-06$ | $4.4 \mathrm{e}-06$ | $5.9 \mathrm{e}-06$ |
| $\mathbf{A}_{\mathbf{1}}(\mathbf{0}, \mathbf{1}, \mathbf{0})$ | n | 100 | 400 | 900 | 1600 | 2500 |
| GFP | IT | 13 | 14 | 15 | 15 | 15 |
| $\omega=1$ | CPU | 0.0045 | 0.0582 | 0.3728 | 5.4135 | 31.8341 |
|  | Res | $5.3 \mathrm{e}-06$ | $7.6 \mathrm{e}-06$ | $9.3 \mathrm{e}-06$ | $2.3 \mathrm{e}-06$ | $2.6 \mathrm{e}-06$ |
| $\mathbf{M G F P}$ | IT | 12 | 13 | 13 | 13 | 13 |
| $\omega=1$ | CPU | 0.0041 | 0.0817 | 0.4127 | 7.2064 | 42.3931 |
| $\alpha=0.1$ | Res | $3.1 \mathrm{e}-06$ | $3.2 \mathrm{e}-06$ | $5.3 \mathrm{e}-06$ | $7.4 \mathrm{e}-06$ | $9.5 \mathrm{e}-06$ |
| $\mathbf{A}_{\mathbf{1}}(\mathbf{1}, \mathbf{1}, \mathbf{1})$ | $\mathbf{n}$ | 100 | 400 | 900 | 1600 | 2500 |
| GFP | IT | 9 | 9 | 10 | 10 | 10 |
| $\omega=1$ | CPU | 0.0046 | 0.0580 | 0.447 | 5.2868 | 31.3160 |
|  | Res | $2.1 \mathrm{e}-06$ | $6.7 \mathrm{e}-06$ | $1.8 \mathrm{e}-06$ | $2.5 \mathrm{e}-06$ | $3.2 \mathrm{e}-06$ |
| $\mathbf{M G F P}$ | IT | 9 | 9 | 9 | 10 | 10 |
| $\omega=1$ | CPU | 0.0050 | 0.0530 | 0.3177 | 5.3091 | 31.5628 |
| $\alpha=0.02$ | Res | $1.6 \mathrm{e}-06$ | $5.0 \mathrm{e}-06$ | $8.3 \mathrm{e}-06$ | $1.8 \mathrm{e}-06$ | $2.3 \mathrm{e}-06$ |
| $\mathbf{A}_{\mathbf{1}}(\mathbf{1}, \mathbf{0}, \mathbf{1})$ | $\mathbf{n}$ | 100 | 400 | 900 | 1600 | 2500 |
| GFP | IT | 9 | 9 | 9 | 10 | 10 |
| $\omega=1.1$ | CPU | 0.0045 | 0.0582 | 0.3728 | 5.4135 | 31.8341 |
|  | Res | $5.3 \mathrm{e}-06$ | $7.6 \mathrm{e}-06$ | $9.3 \mathrm{e}-06$ | $2.3 \mathrm{e}-06$ | $2.6 \mathrm{e}-06$ |
| $\mathbf{M G F P}$ | IT | 9 | 9 | 9 | 9 | 10 |
| $\omega=1.1$ | CPU | 0.0040 | 0.0507 | 0.3671 | 5.4161 | 31.9112 |
| $\alpha=0.05$ | Res | $4.7 \mathrm{e}-06$ | $6.6 \mathrm{e}-06$ | $8.1 \mathrm{e}-06$ | $9.3 \mathrm{e}-06$ | $2.1 \mathrm{e}-06$ |

## 5. CONCLUSION

In this article, we introduced a modified general fixed point method based on new matrix splitting for solving the $\operatorname{LCP}\left(q, A_{1}\right)$ with parameter matrices $\Omega_{1}$ and $\Omega_{2}$. Also, we showed how the iterative form is linked to the new matrix splitting and the parameter matrices $\Omega_{1}$ and $\Omega_{2}$. This iterative form preserves the big and sparse structure of $A_{1}$ during the iteration process. Moreover, we showed the convergence condition for $P$-matrix and presented sufficient convergence domains for $\Omega_{1}$ and $\Omega_{2}$ when system matrix $A_{1}$ is $H_{+}$-matrix. At the end, two examples are discussed to show the efficiency of our proposed method.

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Table 2. Results for GFP [17] and MGFP with $\phi=\alpha\left(L_{1}+U_{1}^{T}\right)$.

| $\mathrm{A}_{1}(\mathbf{1}, \mathbf{1},-1)$ | n | 100 | 400 | 900 | 1600 | 2500 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { GFP } \\ & \omega=1 \end{aligned}$ | IT | 13 | 14 | 15 | 15 | 15 |
|  | CPU | 0.0061 | 0.0807 | 0.5871 | 8.3912 | 48.8538 |
|  | Res | 5.0e-06 | $6.2 \mathrm{e}-06$ | $3.5 \mathrm{e}-06$ | 5.0e-06 | 6.5e-06 |
| MGFP <br> $\omega=1$ <br> $\alpha=0.1$ | IT | 11 | 12 | 13 | 15 | 15 |
|  | CPU | 0.0051 | 0.0607 | 0.5392 | 8.44191 | 49.4862 |
|  | Res | 3.0e-06 | 4.8e-06 | $7.9 \mathrm{e}-06$ | $3.1 \mathrm{e}-06$ | 8.7e-06 |
| $\mathrm{A}_{1}(\mathbf{0}, \mathbf{1 , 0})$ | n | 100 | 400 | 900 | 1600 | 2500 |
| $\begin{aligned} & \text { GFP } \\ & \omega=1 \end{aligned}$ | IT | 10 | 10 | 10 | 11 | 11 |
|  | CPU | 0.0041 | 0.0548 | 0.3621 | 5.999 | 35.2935 |
|  | Res | 1.9e-06 | 5.8e-06 | $9.4 \mathrm{e}-06$ | 2.6e-06 | 3.3e-06 |
| $\begin{aligned} & \text { MGFP } \\ & \omega=1 \\ & \alpha=0.1 \end{aligned}$ | IT | 8 | 9 | 9 | 9 | , |
|  | CPU | 0.0045 | 0.0.0507 | 0.3412 | 4.5795 | 28.4975 |
|  | Res | 7.8e-06 | $2.05 \mathrm{e}-06$ | 2.7e-06 | 3.4e-06 | 4.0e-06 |
| $\mathrm{A}_{1}(\mathbf{1}, \mathbf{1}, \mathbf{1})$ | n | 100 | 400 | 900 | 1600 | 2500 |
| $\begin{aligned} & \text { GFP } \\ & \omega=1 \end{aligned}$ | IT | 7 | 7 | 8 |  |  |
|  | CPU | 0.0041 | 0.0407 | 0.3266 | 4.1892 | 24.6431 |
|  | Res | 2.7e-06 | 6.8e-06 | $9.7 \mathrm{e}-06$ | $1.3 \mathrm{e}-06$ | 1.7e-06 |
| $\begin{aligned} & \text { MGFP } \\ & \omega=1 \\ & \alpha=0.1 \\ & \hline \end{aligned}$ | IT | 6 | 7 | 7 | 7 | 7 |
|  | CPU | 0.0043 | 0.0474 | 0.2796 | 3.4536 | 22.6106 |
|  | Res | 9.7e-06 | $1.3 \mathrm{e}-06$ | $1.6 \mathrm{e}-06$ | $2.0 \mathrm{e}-06$ | 2.2e-06 |
| $\mathrm{A}_{\mathbf{1}}(\mathbf{1 , 0 , 1 )}$ | n | 100 | 400 | 900 | 1600 | 2500 |
| $\begin{aligned} & \text { GFP } \\ & \omega=1.1 \end{aligned}$ | IT | 8 | 8 |  |  |  |
|  | CPU | 0.0041 | 0.0407 | 0.2699 | 4.4741 | 24.6650 |
|  | Res | 1.9e-06 | $2.7 \mathrm{e}-06$ | $3.2 \mathrm{e}-06$ | $3.8 \mathrm{e}-06$ | 4.2e-06 |
| MGFP <br> $\omega=1.1$ <br> $\alpha=0.1$ | IT | 8 | 8 | 8 | 8 | 8 |
|  | CPU | 0.0047 | 0.05303 | 0.2554 | 4.0929 | 24.3851 |
|  | Res | 1.7e-06 | 2.4e-06 | $2.9 \mathrm{e}-06$ | $3.4 \mathrm{e}-06$ | $3.8 \mathrm{e}-06$ |

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