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where U is deduced set contained in  $X_{\ell}(x) = 0$  and  $\nabla f(x^*)$  is the quadrate of f(x) in the point  $x^*$ . Then  $x^*$  is an optimal solution for  $\Gamma$  if and only if  $x^*$  is an optimal solution for  $\Gamma$ ) provided either one of the following

## AN OVERVIEW OF SEPARABLE FRACTIONAL PROGRAMMING PROBLEM

I. M. STANCU-MINASIAN (Control of the control of th

R. Jagannathan (1965b) develops in procedure for solving the

Abstract. This paper represents a sequel to the earlier papers by the author [15 — 18], for presenting the results obtained in fractional programming. Following the classification given by Stancu-Minasian (1981d), in the paper [15] we have presented several applications of fractional programming. We have also presented the methods used for solving: a) the problems of linear fractional programming [16], b) the linear fractional programming problems with several objective functions [17] (1981b)), c) complex fractional programming [18]. In this paper we present the state of the art in separable fractional programming.

Separable programming constitues a linear programming extension for handling certain types of nonlinear functions within the framework of a general linear format. Separable programming problems are nonlinear programming problems with separable objective and constraint functions. First of all, we recall the definition of separability. The func-

tion  $F(x_1, \ldots, x_n)$  is separable if  $F(x_1, \ldots, x_n) = \sum_{j=1}^n F_j(x_j)$  i.e. it can be represented as a sum of functions each involving only one variable in its argument. Then we define the nonlinear separable programming problem.

(1) Maximize (minimize) 
$$F(x) = \sum_{j=1}^{n} F_j(x_j)$$
 subject to

$$x \in D = \left\{x \mid h_i(x) = \sum_{j=1}^n h_{ij}(x_j) \leqslant b_i, \ i = 1, \ldots, m \ ; \ x_j \geqslant 0, \ j = 1, \ldots, n 
ight\}$$

where the functions  $h_i$  are assumed to be separable and continuous.

The basic idea in solving the above problem is to use piecewise linear approximation to each nonlinear function of the objective function and to solve the linear programming problem so obtained. A. Charnes and C. Lemke (1954) proposed an approximation technique for minimizing nonlinear separable convex functionals subject to linear constraints. C. E. Miller (1963) generalized this approach, to include the nonlinear separable functions. Further, we recall a linearization technique due to Kortanek and Evans (1967).

THEOREM 1 (Kortanek and Evans 1967)). Let f be a continuously differentiable function defined on the open convex set  $X \subseteq \mathbb{R}^n$ . Consider the two programs: I maximize  $\{f(x): x \in C\}$  and II maximize  $\{\nabla f(x^*) \cdot x: x \in C\}$ where C is a closed set contained in X,  $x^* \in C$  and  $\nabla f(x^*)$  is the gradient of f(x) in the point  $x^*$ . Then  $x^*$  is an optimal solution for I if and only if x\* is an optimal solution for II provided either one of the following conditions holds: a) f is pseudoconcave on X: b) f is quasiconcave on X and  $\nabla f(x^*) (\not\equiv 0.) \text{AHA}$  FIRACOP VALITATION N.

Recently, these methods were generalized for solving separable fractional programming problems (see Almogy and Levin (1971), Anand and Swarup (1970), Arora and Aggarwal (1977), Černov and Lange (1970), Černov (1971, 1972) Gogia (1969), Jagannathan (1965a, b), Kaul and Datta (1981), Tigan (1977, 1981) and Wadhwa (1969)).

R. Jagannathan (1965b) develops a procedure for solving the

problem

2) Minimize  $\{F(x_1,...,x_n) = \sum_{i=1}^{n} (c_i/(x_i + s_i)) \ (c_i > 0, s_i > 0) : A, x \le bx \ge 0\}$ ding. Following the classification given by Stonen Minasian (1981d), in the He demonstrates the following theorem: s between and ow [6] raged

THEOREM 2 (Jagannathan (1965b)). If  $F(\bar{x}) = \sum_{j=1}^{n} (c_j/(x_j + s_j))$  and

$$Q^{0}(\bar{x}) := \sum_{j=1}^{n} \lambda_{j} x_{j}^{2} - \sum_{j=1}^{n} p_{j} x_{j} \text{ where } \lambda_{j} = -\frac{1}{s_{j}} \frac{\partial F_{0}}{\partial x_{j}}, \ p_{j} = \frac{-(2x_{j} + s_{j})}{s_{j}} \frac{\partial F}{\partial x_{j}^{0}} \text{ and } s_{j}$$

noisetely minimizer 
$$q$$
 so the second  $\bar{x}^{0\prime}:=(x_1^0,\dots,x_n^0)$ , then  $(i)$   $F(\bar{x}^0):=\sum\limits_{j=1}^n (c_j/s_j)+Q^0(\bar{x}^0)$  multiplies to the second  $\bar{x}^{0\prime}:=(x_1^0,\dots,x_n^0)$  and second  $\bar{x}^{0\prime}:=(x_1^0,\dots,x_n^0)$ 

$$(ii) \ \ F(ar{x}) \leqslant \sum\limits_{j=1}^{n} (c_j/s_j) + Q^0(ar{x}) for \ all \ ar{x} \geqslant 0, (iii) rac{\partial Q^0(ar{x})}{\partial ar{x}^0} = rac{\partial F(ar{x})}{\partial ar{x}^0} \ .$$

According to this theorem, R. Jagannathan reduces the solving of separable fractional programming problem (2) to a sequence of convex quadratic programming problems. Minimize  $\{Q^i(x): Ax \leq b; x \geq 0\}$  where the solution of the  $k^{\text{th}}$  problem is used to define  $Q^{(k+1)}(x)$  for the  $(k+1)^{\text{th}}$ problem. The algorithm stops when the reduction in the successive values of Q(x) becomes very small. We remark that due to the function  $F(x_1,...,x_n)$ is convex we can apply the method of A. Charnes and O. Lemke (1954) to solve the problem (1).

P. Anand and K. Swarup (1970) consider the following problem:

(3) Maximize 
$$\left\{ F(x_1, \dots, x_n) = \sum_{j=1}^n [(x_j + \alpha_j)/(x_j + \beta_j)], (\alpha_j > 0, \beta_j > 0) \right\}$$

subject to  $x \in \{x \mid Ax \leqslant b, x \geqslant 0\}$ 

Here the objective function is neither concave nor convex. In order to solve this problem, they reduce it to a sequence of concave quadratic programming problems. The sequence of points solving the quadratic problems is shown to converge to a local solution of the problem (3).

The relation between the problem (3) and the concave quadratic

problem is given by the following theorem.

THEOREM 3 (Anand and Swarup (1970)). If 
$$F(x) = \sum_{j=1}^{n} \frac{x_j + \dot{\alpha}_j}{x_j + \dot{\beta}_j}$$
 and  $Q^0(x) = -\sum_{j=1}^{n} \lambda_j x_j^2 + \sum_{j=1}^{n} [x_j(p_j - \alpha_j \lambda_j) + p_j \alpha_j]$  where  $\lambda_j = -\frac{1}{\alpha_j - \dot{\beta}_j} \frac{\partial F}{\partial x_j^0}$  and  $p_j = -\frac{(2x_j^0 + \beta_j)}{\alpha_j - \beta_j} \frac{\partial F}{\partial x_j^0}$ , then (i)  $F(x) \geqslant Q^0(x)$ , (ii)  $F(x) = Q^0(x^0)$ , (iii)  $\nabla F(x^0) = \nabla Q^0(x^0)$ .

N.K. Gogia (1969) considers the following problem

(4) Maximize  $F(x_1, \ldots, x_n) = \left(\sum_{j=1}^n f_j(x_j)\right) / \left(\sum_{j=1}^n g_j(x_j)\right)$  subject to  $x \in D$  where  $D=\left\{x\in R^n\mid \sum\limits_{j=1}^nh_{ij}(x_j)\leqslant b_i,\;i=1,\ldots,m, |x_j|\geqslant 0, |j|=1,\ldots,n
ight\}$  and the problem functions are assumed continuous. Each of the functions  $f_j$ ,  $g_j$  and  $h_{ij}$  is approximated by a set of piecewise linear functions following C. E. Miller (1963), and the resulting linear fractional program is

(5) Maximize  $E^* = \left(\sum_{j=1}^n \sum_{k=0}^{r_j} \lambda_{kj} f_{kj}\right) / \left(\sum_{j=1}^n \sum_{k=0}^{r_j} \lambda_{kj} g_{kj}\right) \text{subject to } \sum_{j=1}^n \sum_{k=0}^{r_j} \lambda_{kj} h_{kij} \leqslant$  $\leqslant b_i,\ i=1,\ldots,m$ ;  $\sum_{k,j} \lambda_{kj}=1, j=1,\ldots,n$ ;  $\lambda_{kj}\geqslant 0$  for all k,j. Solving this problem by applying simplex algorithm, with restricted basis entry in the basis such that not more than two \(\lambda\_{ks}\) are positive we obtain the solu-

tion  $\lambda_{kj}$ . If the two  $\lambda_{ks}$  are positive then they must be adjacent. Then  $x_j = \sum_{k=0}^{j} \lambda_{kj} x_{kj}$  will be the approximate solution of the original problem.

Ju. P. Černov and E. G. Lange (1970) give an approximative method for solving a transport problem with separable fractional function

(6) Minimize 
$$F(x) = \left(\sum_{i=1}^{m} \sum_{j=1}^{n} f_{ij}(x_{ij})\right) / \left(\sum_{i=1}^{m} \sum_{j=1}^{n} g_{ij}(x_{ij})\right)$$
 where  $x = (x_{ij})$  and  $x \in D = \left\{x \in R \middle| \sum_{j=1}^{n} x_{ij} = a_i, \sum_{i=1}^{m} x_{ij} = b_i\right\}$ ,  $R = \left\{x \mid a_{ij} \leqslant x_{ij} \leqslant \beta_{ij}\right\}$ . The functions  $f_{ij}$ ,  $g_{ij}$  are linearized using the Alform (Hodlar (1004)).

The functions  $f_{ij}$ ,  $g_{ij}$  are linearized using the  $\delta$ -form (Hadley (1964)) and the new problem has the form and my loor as guidant to saill o

 $\begin{array}{ll} \text{Minimize} & \{(\sum\sum c'_{rs}y_{rs}+c'_{\mathbf{0}})/(\sum\sum c'_{rs}y_{rs}+c''_{\mathbf{0}}) \mid \sum_{s} y_{rs}=A_{r}, \sum_{r}y_{rs}=B_{s}, \\ y_{rs}\geqslant 0\}. \end{array}$ 

and  $y_{rs}$  must satisfy a certain supplimentary condition.

Ju. P. Černov (1971) considers the following separable problem

(7) Minimize 
$$F(x_1, ..., x_n) = \left(\sum_{j=1}^n f_j(x_j)\right) / \left(\sum_{j=1}^n g_j(x_j)\right)$$
  
subject to  $x \in \left\{x \mid \sum_{j=1}^n h_{ij}(x_j) \leqslant b_i, i = 1, ..., n ; \alpha_j \leqslant x_j \leqslant \beta_j, j = 1, ..., n\right\}$ .

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He utilizes the same method, namely to approximate all separable functions by linear functions with a greater number of variables, and by means of it reduces the problem to linear fractional programming problem.

Vijay Wadhwa (1969) considers a class of mathematical programming problems where the objective function is the sum of separable quadraticlinear fractional functionals and the set of constraints is a convex 

(8) Minimize 
$$\left\{F(x_1,\ldots,x_n) = \sum_{j=1}^n [(k_j + m_j x_j^2)/(x_j + c_j)] \mid Ax = b \; ; \; x \geqslant 0 \right\}$$

It is supposed that  $c_i > 0$ ,  $k_i > 0$ ,  $m_i \ge 0$ , which make the above function convex. The method proposed requires the solution of a sequence of quadratic programming problems and is based on the following theorem.

THEOREM 4 (Wadhwa (1969)). If 
$$F(x) = \sum_{j=1}^{n} [(k_j + m_j x_j^2)/(x_j + c_j)],$$

where 
$$k_j$$
,  $c_j > 0$  and  $m_j \ge 0$  are constants and  $G_0(x) = \sum_{j=1}^n \alpha_j x_j^2 = \sum_{j=1}^n \beta_j x_j$ ,

where 
$$\alpha_j = [m_j - \nabla F(x_j^0)]/c_j$$
,  $\beta_j = [2x_j^0(k_j + m_j c_j^2)/c_j(x_j^0 + c_j)^2] - \nabla F(x_j^0)$ , and  $x^0 = (x_1^0, \dots, x_n^0) \ge 0$ , then a)  $F(x) \leqslant G_0(x) + \sum_{j=1}^n (k_j/c_j)$ , for all  $x \ge 0$ ,

$$F(x^0) = G_0(x^0) + \sum_{j=1}^n (k_j/c_j), \ c) \ \ \nabla F(x^0) = \nabla G_0(x^0).$$

Y. Almogy and O. Levin (1971) consider the following separable problem (9) Maximize  $\left\{F(x_1,\ldots,x_n) = \sum_{i=1}^n \left[\frac{(c_i'x+\alpha_i)}{(d_i'x+\beta_i)}\right] \mid Ax \leqslant b, \ x \geqslant 0\right\}$ where x is an n-component column vector,  $c_i$  and  $d_i$  are n-component column vectors (i = 1, ..., k), A is an  $m \times n$  matrix, b is an m-component column vector, and  $\alpha_i$ ,  $\beta_i$  are given constants. The property of the solution of the solu

They consider the value of each function  $f_i(x_i)$  as a parameter  $f_i$ that can be varied, irrespective of whether or not it is attainable on feasible set. In this manner, the problem can be transformed into equivalent ones of maximizing multiparameter linear or concave functions subject to additional feasibility constraints. If the linear fraction's denominators are restricted to nonnegative coefficients, the problems can be transformed into those of finding a root (in the case of separability this root is unique) of a monotone-decreasing convex parametric function. When the number of terms in the objective function is equal or less than three. Y. Almogy and O. Levin give an efficient algorithm. In the special case when the coefficient matrix A consists of nonnegative elements (as in most transportation problems), Y. Almogy and O. Levin utilized the combinatorial property of the parametric presentation in outling a method of solution. [1] Align J. W. M. De James A. T. Alexandron, Phys. Rev. B 98 (1994) 115 (1994).

St. Tigan (1977) comes to a separable fractional programming problem related to the resources dynamic distribution problem, wherein the resources utilization efficiency (mean income to the time unit) is maxiFRACTIONAL PROGRAMMING

mized. The model is the following:

Maximize 
$$F(x_1, \ldots, x_n) = \left(\sum_{i=1}^n f_i(x_i) + \alpha\right) / \left(\sum_{i=1}^n g_i(x_i) + \beta\right)$$
  
subject to  $\sum_{i=1}^n w_i x_i \leqslant b$ ,  $\alpha_i \geqslant x_i \geqslant \beta_i \geqslant 0$ ,  $x_i$  integers where  $w_i$ ,  $\alpha_i$ ,  $\beta_i$   $(i = 1)$ 

= 1,...,n),  $\alpha$ ,  $\beta$  and b are given real numbers. He gives an algorithm for solving this problem which consists in solving a finite number of

ordinary dynamic programming problems.

Also, I mention a recently paper of Savita Arora and S. P. Aggarwal (1977) dealing with maximizing the sum of a finite number of separable linear concave (convex) fractional Functions. The nonconvex piecewise linear separable programming problem is the following:

Maximize 
$$\left\{V_s = \sum_{j=1}^s \frac{\gamma_j}{c_j' X_j + \alpha_j} \left| \sum_{j=1}^s \bar{B}_j X_j = \bar{b}, \; B_j X_j = b_j, \; j=1,\ldots,s, X_j \right. \right\}$$
 where  $c_j,\; X_j$  are each  $n_j \times 1$  component vectors,  $\bar{b}$ 

and  $b_i$  are vectors with m and  $m_i$  components,  $\overline{B}_i$  and  $B_i$  are  $m \times n_i$ and  $m_i \times n_i$  matrices,  $\gamma_i$  and  $\alpha_i$  are scalar constants. A decomposition principle is derived with the dynamic programming approach to solving this problem. This results in a series of parametric quadratic fractional subprogrammes whose recursive solution yields the solution to the original

Finally, I give a recent paper of R. N. Kaul and Neelam Datta (1981) which considers the following nonlinear fractional programming 

Maximize 
$$F(x) = \left(\sum_{j=1}^{n_1} f_j(x_j) + \sum_{j=1}^{n_2} h_j(y_j)\right) / \sum_{j=1}^{n_1} g_j(x_j)$$
 subject to  $\sum_{j=1}^{n_1} e_{ij}(x_j) + \sum_{j=1}^{n_2} e_{ij}(x_j)$ 

 $+\sum_{j=1}^{n_2}a_{ij}y_j\leqslant b_i,\,\,i=1,\ldots,m\,;\,\,x_j\geqslant 0,\,\,\,j=1,\ldots,n_1\,;\,\,y_j\geqslant 0,\,\,\,j=1,\ldots,n_2$ where the functions  $f_i(x_i)$   $(j = 1, ..., n_1)$  are concave and  $h_i(y_i)$   $(j = 1, ..., n_2)$ 

 $g_i(x_i)$   $(j = 1, ..., n_i)$  and  $e_{ii}$   $(i = 1, ..., m; j = 1, ..., n_i)$  are convex and  $g_i(x_i)$  are assumed to be positive over the feasibility region.

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