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# SOME PROCEDURES FOR SOLVING SPECIAL MAX-MIN FRACTIONAL RANK-TWO REVERSE-CONVEX PROGRAMMING PROBLEMS

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Dedicated to Professor Elena Popoviciu on the occasion of her 80th birthday

**Abstract.** In this paper, we suggest some procedures for solving two special classes of max-min fractional reverse-convex programs. We show that a special bilinear fractional max-min reverse-convex program can be solved by a linear reverse-convex programming problem.

For a linear fractional max-min reverse-convex program, possessing two reverse-convex sets, we propose a parametrical method. The particularity of this procedure is the fact that the max-min optimal solution of the original problem is obtained by solving at each iteration two linear reverse-convex programs with a rank-two monotonicity property.

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**Keywords.** Reverse-convex programming, max-min programming, bilinear fractional programs, linear fractional max-min programs.

### 1. INTRODUCTION

In this paper, we propose some methods for solving two special classes of max-min fractional reverse-convex programs.

Let  $X \subseteq \mathbb{R}^n, Y \subseteq \mathbb{R}^m$  be given convex sets. Let T be a reverse-convex set in  $\mathbb{R}^n$  (i.e. the complement of a convex set in  $\mathbb{R}^n$ ) and S be a reverse-convex set in  $\mathbb{R}^m$  and  $h: X \times Y \to \mathbb{R}$ . The general max-min reverse-convex programming consists in finding an optimal max-min solution for:

**PRC**. 
$$\max_{x \in X \cap T} \quad \min_{y \in Y \cap S} h(x, y).$$

We recall (see, e.g. [14, 17]) that a point  $(x', y') \in (X \cap T) \times Y$  is said to be an *optimal* max-min *solution* for **PRC** problem, if the following two conditions hold:

(i) 
$$h(x', y') = \max_{x \in X \cap T} \min_{y \in Y \cap S} h(x, y)$$
  
(ii) 
$$\min_{y \in Y \cap S} h(x', y) = h(x', y').$$

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In the particular case, when **PRC** is a simple maximization problem only (i.e. Y has a single element, X is a polyhedral set and T is a reverse-convex set defined by a convex or quasi-convex constraint which has a rank-two monotonicity property), we recall that procedures for solving the maximum (or minimum) reverse-convex programming problems were given, for instance, in the papers [4], [7], [8], [10], [18], [6]. Duality aspects of the reverse-convex programming can be found in the paper by Penot [9]. For the max-min reverseconvex programming we mention the refs. [5], [6].

In section 2, for a special fractional type of the objective function h and for some particular cases of the sets X, Y, T and S, we obtain a particular case of problem **PRC** of bilinear fractional form, for which we propose, in Section 3, a method of finding optimal max-min solutions by reducing these problems to the particular case when X is a polyhedral set and T is a reverseconvex set defined by a convex or quasi-convex constraint that has a rank-two monotonicity property. These auxiliary problems can be solved by Kuno-Yamamoto procedure [8].

In section 4, we consider a max-min linear fractional reverse-convex programming problem, having two reverse-convex sets T and S, and for which we propose a parametric procedure.

Some concluding remarks are made in the last section.

### 2. BILINEAR FRACTIONAL MAX-MIN REVERSE-CONVEX PROGRAM PFM

Next we consider the following bilinear fractional max-min reverse-convex program (see, [6]) in which the reverse set  $S = \mathbb{R}^m$ :

**PFM.** Find

(3)

(1) 
$$V = \max_{x \in X \cap T} \min_{y \in Y} \left( \frac{y^{\mathrm{T}} C x + c^{\mathrm{T}} x + e^{\mathrm{T}} y + e_{0}}{\max\{\alpha_{i}^{\mathrm{T}} y + \beta_{i} | i \in I\}} \right)$$

where  $I = \{1, 2, ..., r\}, C \in \mathbb{R}^{m \times n}, c \in \mathbb{R}^n, e, \alpha_i \in \mathbb{R}^m \text{ and } e_0, \beta_i \in \mathbb{R}, (i \in I).$ The sets X and Y are defined by

(2) 
$$X = \{x \in \mathbb{R}^n | Ax = a, x \ge 0\},$$

$$Y = \{ y \in \mathbb{R}^m | Dy \ge d, y \ge 0 \}$$

where  $A \in \mathbb{R}^{s \times n}, a \in \mathbb{R}^s, D \in \mathbb{R}^{q \times m}, d \in \mathbb{R}^q$ .

The reverse-convex set

(4) 
$$T = \{ x \in X^0 | f(x) \le 0 \},\$$

is defined by a function  $f : \mathbb{R}^n \to \mathbb{R}$ , which is continuous, strictly quasiconcave and has rank-two monotonicity on an open convex set  $X^0 \subseteq \mathbb{R}^n$ , which includes the set X.

We recall that f is said to be strictly quasiconcave on  $X^0$  if for each  $x, y \in X^0$  with  $f(x) \neq f(y)$  we have  $f((1-t)x + ty) > \min\{f(x), f(y)\}$ , for any  $t \in (0, 1)$ .

DEFINITION 1. [8] The function f possess a rank-two monotonicity on  $X^0$ with respect to linearly independent vectors  $\lambda_1, \lambda_2 \in \mathbb{R}^n$ , if for any points  $x', x'' \in X^0$ , the inequality  $\lambda_i x' \leq \lambda_i x''$ , i = 1, 2 implies  $f(x') \leq f(x'')$ .

We have the following representation for a function f possessing a rank-two monotonicity on  $X^0$  with respect to linearly independent vectors  $\lambda_1, \lambda_2 \in \mathbb{R}^n$ .

LEMMA 2. [8] If the function f possess a rank-two monotonicity on  $X^0$ with respect to linearly independent vectors  $\lambda_1, \lambda_2 \in \mathbb{R}^n$ , then there exists a function  $g : \mathbb{R}^2 \to \mathbb{R}$ , which is continuous and strictly quasiconcave on  $\Gamma^0 = \{(\lambda_1 x, \lambda_2 x) | x \in X^0\}$  and satisfies the following two conditions: (i)  $f(x) = g(\lambda_1 x, \lambda_2 x)$ , for  $x \in X^0$ , (ii)  $g(\theta) \leq g(\eta)$  if  $\theta, \eta \in \Gamma^0$  and  $\theta \leq \eta$ .

Concerning the problem **PFM**, we make the following assumptions:

A1. Y is a bounded non-empty set,

A2. the function f is continuous, strictly quasi-concave and has rank-two monotonicity on the open convex set  $X^0$ ,

A3.  $\max\{\alpha_i y + \beta_i | i \in I\} > 0, \forall y \in Y.$ 

# 3. LINEAR REVERSE-CONVEX PROGRAMMING APPROACH

We proposed in [6] a procedure for solving problem **PFM** based on the Charnes-Cooper [1] variable change.

Thus, if we perform in the problem (1)–(4), the Charnes-Cooper variable change v = ty,  $t \ge 0$ ,  $v \in \mathbb{R}^s$ , (see, also Schaible [11], Stancu-Minasian [12] and Tigan [14], [16]) it follows by assumptions A1 and A3 that problem **PFM** is equivalent with:

**PA**. Find

(5) 
$$V' = \max_{x \in X \cap T(v,t) \in Y'} \left( v^{\mathrm{T}} C x + c^{\mathrm{T}} x t + e^{\mathrm{T}} v + e_0 t \right),$$

where the set Y' is defined by:

$$(6) Dv - dt \ge 0,$$

(7) 
$$\max\{\alpha_i^{\mathrm{T}}v + \beta_i t | i \in I\} \ge 1,$$

$$(8) v \ge 0, t \ge 0.$$

We can show, by assumptions A1 and A3, that V = V' and that for any optimal max-min solution  $(x^*, y^*)$  of problem **PFM** there exists an optimal max-min solution  $(x^*, v^*, t^*)$  of problem **PA** such that  $y^* = \frac{v^*}{t^*}$  and conversely.

By Golstein [3] (see, also Ţigan and Stancu-Minasian [17]), the constraint (7) in Problem **PA** can be rewritten as the following maximum bilinear constraint:

(9) 
$$\max_{\theta \in Z} \quad \sum_{i=1}^{r} \theta_i (\alpha_i v + \beta_i t - 1) \ge 0,$$

where  $Z = \{ \theta \in \mathbb{R}^r | \sum_{i=1}^r \theta_i = 1, \theta_i \ge 0, i = 1, 2, ..., r \}.$ Let us denote

(10) 
$$\psi(v,t) = \max_{\theta \in Z} \quad \sum_{i=1}^{r} \theta_i (\alpha_i^{\mathrm{T}} v + \beta_i t - 1), \forall (v,t) \in Y'.$$

By linear programming duality, for any  $(v, t) \in Y'$ , we have

(11) 
$$\psi(v,t) = \min w$$

subject to

(12) 
$$w \ge \alpha_i^{\mathrm{T}} v + \beta_i t - 1, \forall i \in I, w \ge 0.$$

From (10)-(12), it follows that the inequation (9) can be expressed by the following system:

$$w \ge 0, w \ge \alpha_i^{\mathrm{T}} v + \beta_i t - 1, \forall i \in I.$$

Therefore, problem **PA** can be reduced to the following max-min bilinear reverse-convex program

PMM1. Find

(13) 
$$V' = \max_{x \in X \cap T} \quad \min_{(v,t,w) \in Y''} \left( v^{\mathrm{T}}(Cx+e) + t(c^{\mathrm{T}}x+e_0) \right),$$

where the set Y'' is defined by:

$$(14) Dv - dt \ge 0,$$

(15) 
$$w \ge \alpha_i^{\mathrm{T}} v + \beta_i t - 1, \forall i \in I$$

(16) 
$$v \ge 0, t \ge 0, w \ge 0.$$

From (13)–(16), by linear programming duality, for any  $x \in X \cap T$ , problem **PMM1** can be transformed into the following linear reverse-convex program: PML. Find

(17) 
$$V' = \max_{x,u,z} \left( -z_1 - z_2 - \dots - z_r \right),$$

subject to

$$u^{\mathrm{T}}D - z^{\mathrm{T}}\Omega \leq Cx + e,$$
  
$$-u^{\mathrm{T}}d - z^{\mathrm{T}}\Lambda \leq c^{\mathrm{T}}x + e_{0}, i \in I$$
  
$$x \in X \cap T,$$
  
$$u \geq 0, z \geq 0, u \in \mathbb{R}^{q}, z \in \mathbb{R}^{r}$$

In problem **PML**, we denote by  $\Omega \in \mathbb{R}^{r \times s}$  the matrix having the rows  $\alpha_i$  $(i \in I)$ , and by  $\Lambda$  the vector  $\Lambda = (\beta_1, ..., \beta_r)^{\mathrm{T}} \in \mathbb{R}^r$ .

Therefore, we proved the following theorem:

THEOREM 3. If the problem **PFM** satisfies the assumptions A1-A3, then problem **PFM** can be solved by solving the linear reverse-convex program with a rank-two monotonicity **PML**.

In order to solve problem **PFM**, a procedure similar to Algorithm 1 can be used, by replacing in step1 the problem **PLC** by the linear reverse-convex programming problem **PML**.

ALGORITHM 1. Step 1. Solve the linear reverse-convex programs with a rank-two monotonicity **PML**.

If  $V < \infty$  and the feasible set X' of **PML** is non-empty, let  $x^*$  be the corresponding component of an optimal solution of problem **PML**.

If  $X = \emptyset$ , then take  $V = -\infty$ .

Step 2. i) If  $-\infty < V < \infty$ , then  $(x^*, y^*)$  is an optimal solution of maxmin problem **PFM**, where  $y^*$  is an optimal solution of the generalized linear-fractional program **PFA**.

ii) If  $V = -\infty$ , then **PFM** is unfeasible.

iii) If  $V = \infty$ , then **PFM** has an infinite optimum.

We make the remark that auxiliary linear reverse-convex programming problem **PML** in Algorithm 1 is simpler than the auxiliary linear reverseconvex programming problem in the algorithm proposed for this problem in [6]. Indeed, problem **PML** has only s+r linear constraints the auxiliary problem in [6] posses (s+2)r linear constraints. Therefore, Algorithm 1 seems to be more efficient than algorithm proposed for this problem in [6].

# 4. MAX-MIN LINEAR FRACTIONAL REVERSE-CONVEX PROGRAMMING WITH TWO SEPARATE REVERSE-CONVEX FEASIBLE SETS

In this section, we consider the following max-min linear fractional program **GLF**. Find

$$\max_{x \in X \cap T} \quad \min_{y \in Y \cap S} h(x, y)$$

where

$$h(x,y) = \frac{\alpha^{\mathrm{T}} x + \beta^{\mathrm{T}} y + \beta_0}{\gamma^{\mathrm{T}} x + \eta^{\mathrm{T}} y + \eta_0},$$

verifying the condition

$$\forall^{\mathrm{T}} x + \eta^{\mathrm{T}} y + \eta_0 > 0, \ \forall x \in X, \ \forall y \in Y.$$

In problem **GLF**, X and S are defined by

- (18)  $X = \{ x \in \mathbb{R}^n | Ax = a, x \ge 0 \},$
- (19)  $Y = \{ y \in \mathbb{R}^m | By = b, y \ge 0 \},$

where  $A \in \mathbb{R}^{s \times n}, B \in \mathbb{R}^{p \times m}, a \in \mathbb{R}^s, b \in \mathbb{R}^p, \alpha, \gamma \in \mathbb{R}^n, \beta, \eta \in \mathbb{R}^m, \beta_0, \eta_0 \in \mathbb{R}$ , are given matrices, vectors and real constants respectively.

The reverse-convex set

(20) 
$$T = \{x \in X^0 | f(x) \le 0\}$$

is defined by a function  $f : \mathbb{R}^n \to \mathbb{R}$ , which is continuous, strictly quasiconcave and has rank-two monotonicity on an open convex set  $X^0 \subseteq \mathbb{R}^n$ , which includes the set X and the reverse-convex set

(21) 
$$S = \{ y \in Y^0 | f_1(y) \le 0 \},$$

is defined by a function  $f_1 : \mathbb{R}^m \to \mathbb{R}$ , which is continuous, strictly quasiconcave and has rank-two monotonicity on an open convex set  $Y^0 \subseteq \mathbb{R}^m$ , which includes the set Y.

For solving problem **GLF** we can use a parametric procedure (see, [2], [13], [15]), by which an approximate optimal solution could be found by solving a sequence of the auxiliary reverse-convex programs each of them having only one reverse-convex constraint.

ALGORITHM 2. Let  $\varepsilon > 0$  be a given positive real number, representing a level of approximation to be attain by algorithm.

1. Find a point  $x^0 \in X \cap T$  and a point  $y^0 \in Y \cap S$  and set k := 0.

2. Take

$$t_k = h(x^k, y^k).$$

3. Find

(22) 
$$F(t_k) = \max_{x \in X \cap T} \min_{y \in Y \cap S} [(\alpha - t_k \gamma)x + (\beta - t_k \eta)y + \beta_0 - t_k \eta_0].$$

But the max-min program (22) can be transformed into the following two linear reverse-convex programs

**PL1**. Find

(23) 
$$\max(\alpha - t_k \gamma) x$$

subject to

$$(24) x \in X \cap T$$

PL2. Find

(25) 
$$\min_{y} [(\beta - t_k \eta)y + \beta_0 - t_k \eta_0]$$

subject to

$$(26) y \in Y \cap S.$$

Let  $x^{k+1}, y^{k+1}$  be an optimal solution of the linear reverse-convex program (23)–(24) and (25)–(26), respectively. Obviously, we have  $F(t_k) = (\alpha - t_k \gamma) x^{k+1} + (\beta - t_k \eta) y^{k+1} + \beta_0 - t_k \eta_0$ . 4. i) If  $F(t_k) \leq \varepsilon$ , then  $(x^{k+1}, y^{k+1})$  is an approximate optimal solution of

4. i) If  $F(t_k) \leq \varepsilon$ , then  $(x^{k+1}, y^{k+1})$  is an approximate optimal solution of problem **GLF**.

ii) If  $F(t_k) > \varepsilon$ , then take k:=k+1 and go to the step 2.

#### 5. CONCLUSIONS

In this paper we consider two fractional max-min reverse-convex programming problems.

Firstly, we give a new procedure for solving a particular class of maxmin bilinear fractional reverse-convex programming problems. The particularity of this procedure is the fact that the max-min optimal solution of the original problem is obtained by solving a single linear reverse-convex program with a rank-two monotonicity with an algorithm proposed by Kuno and Yamamoto [8].

Secondly, we consider a parametric procedure for solving a particular class of max-min linear fractional reverse-convex programming problems, possessing two reverse-convex feasible sets. The particularity of this procedure is the fact that the max-min optimal solution of the original problem is obtained by solving at each iteration two linear reverse-convex programs with a rank-two monotonicity with the algorithm of Kuno and Yamamoto.

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